Systems Programming

Threads

Byoungyoung Lee Seoul National University

byoungyoung@snu.ac.kr

https://lifeasageek.github.io

Today

- Concurrent Programming is Hard
- Threads

Concurrent Programming is Hard!

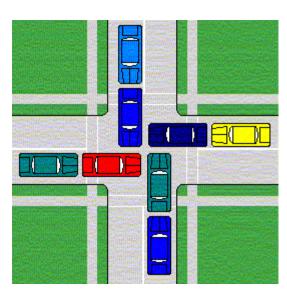
- The human mind tends to be sequential
- The notion of time is often misleading
- Thinking about all possible sequences of events in a computer system is at least error prone and often impossible
- Often leads to concurrency bugs
 - Data race, deadlock, livelock, and starvation

Data Race

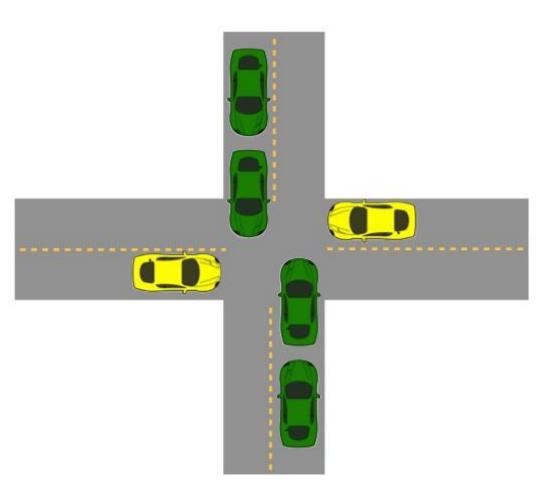


Deadlock





Starvation



- Yellow must yield to green
- Continuous stream of green cars
- Overall system
 makes progress, but
 some individuals
 wait indefinitely

Concurrent Programming is Hard!

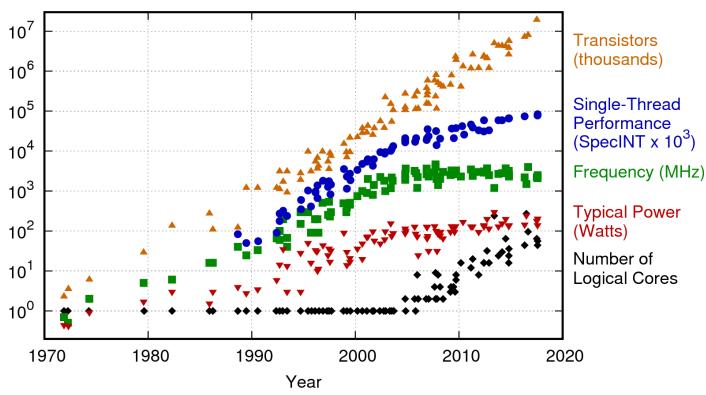
- Classical problem classes of concurrent programs:
 - Races: outcome depends on arbitrary scheduling decisions
 - Example: who gets the last seat on the airplane?
 - Deadlock: improper resource allocation prevents forward progress
 - Example: traffic gridlock
 - Starvation: external events and/or system scheduling decisions can prevent sub-task progress
 - Example: people always jump in front of you in line
- Many aspects of concurrent programming are beyond the scope of our course..
 - We'll cover some of these aspects in the next few lectures.

Concurrent Programming is Hard!

It may be hard, but ...

it is useful and become more and more necessary!





Original data up to the year 2010 collected and plotted by M. Horowitz, F. Labonte, O. Shacham, K. Olukotun, L. Hammond, and C. Batten New plot and data collected for 2010-2017 by K. Rupp

Today

- **■** Concurrent Programming is Hard
- Threads

Process is defined by ...

Process = process context + code/data/stack+ kernel data structures

Code/Data/Stack Stack SP **Shared libraries Process context: Data registers** brk **Condition codes Run-time heap** Stack pointer (SP) Read/write data **Program counter (PC)** PC Read-only code/data **Kernel data structures:** VM structures **Descriptor table**

A Process With Multiple Threads

- Multiple threads can be associated with a process
 - Each thread has its own
 - logical control flow
 - stack (but not protected from other threads)
 - thread id (TID)
 - Threads share the same
 - code, data, and kernel context

Thread 1

stack₁

Thread 1 context:

Data registers₁

Condition codes₁

SP₁

PC₁

Thread 2

stack₂

Thread 2 context:

Data registers₂

Condition codes₂

SP₂

PC₂

Code/Data

shared libraries

run-time heap read/write data

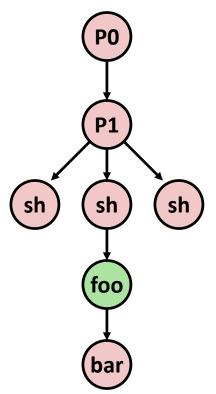
read-only code/data

Kernel data structures:
VM structures
Descriptor table

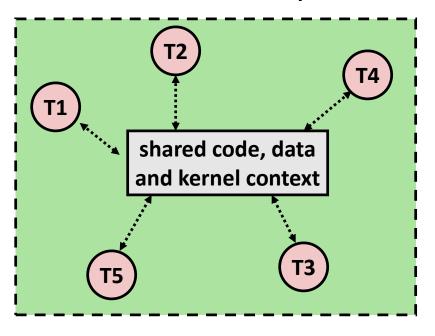
Logical View of Threads

- Threads associated with process form a pool of peers
 - Unlike processes which form a tree hierarchy

Process hierarchy



Threads associated with the process foo

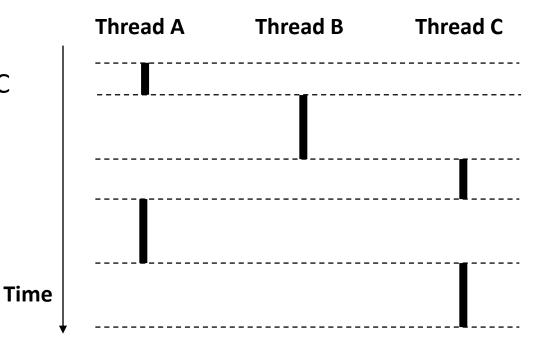


Concurrent Threads

- Two threads are concurrent if their flows overlap in time
- Otherwise, they are sequential

Examples:

- Concurrent: A & B, A&C
- Sequential: B & C



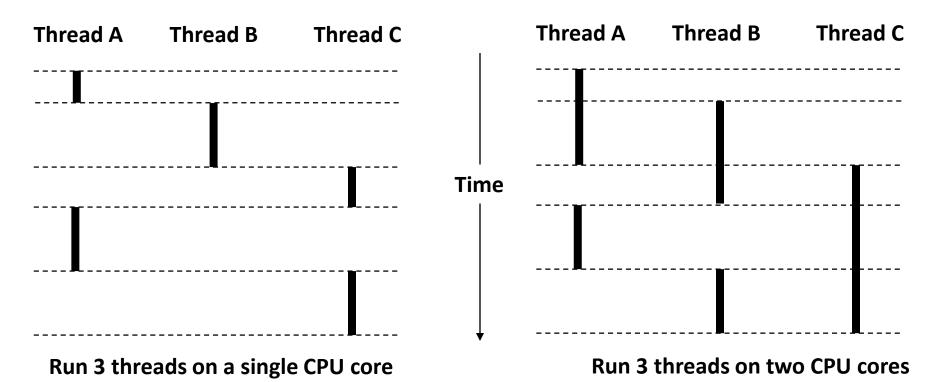
Concurrent Thread Execution

Single Core Processor

Simulate parallelism by time slicing

Multi-Core Processor

Can have true parallelism



Threads vs. Processes

How threads and processes are similar

- Each has its own logical control flow
- Each can run concurrently with others (possibly on different cores)
- Each is context switched

How threads and processes are different

- Threads share all code and data (except local stacks)
 - Processes do not
- Threads are somewhat less expensive than processes
 - Process control (creating and reaping) twice as expensive as thread control
 - Linux numbers:
 - ~20K cycles to create and reap a process
 - ~10K cycles (or less) to create and reap a thread

Posix Threads (pthreads) Interface

- pthreads: Standard interface for ~60 functions that manipulate threads from C programs
 - Creating and reaping threads
 - pthread_create()
 - pthread_join()
 - Determining your thread ID
 - pthread_self()
 - Terminating threads
 - pthread_cancel()
 - pthread_exit()
 - exit() [terminates all threads]
 - return [terminates current thread]
 - Synchronizing access to shared variables
 - pthread_mutex_init
 - pthread_mutex_[un]lock

The pthreads "hello, world" Program

```
* hello.c - pthreads "hello, world" program
 */
                                                       Thread attributes
                                     Thread ID
#include "csapp.h"
                                                        (usually NULL)
void *thread(void *varqp);
int main (int argc, char** argv)
                                                        Thread routine
    pthread t tid;
    pthread create (&tid, NULL, thread, NULL);
    pthread join(tid, NULL);
                                                          Thread arguments
    return 0;
                                                              (void *p)
                                                 hello.c
                                                       Return value
                                                        (void **p)
void *thread(void *varqp) /* thread routine */
    printf("Hello, world!\n");
    return NULL;
                                                 hello.c
```

Execution of Threaded "hello, world"

Main thread

