

Systems Programming

Machine-Level Programming II: Control

Byoungyoung Lee

Seoul National University

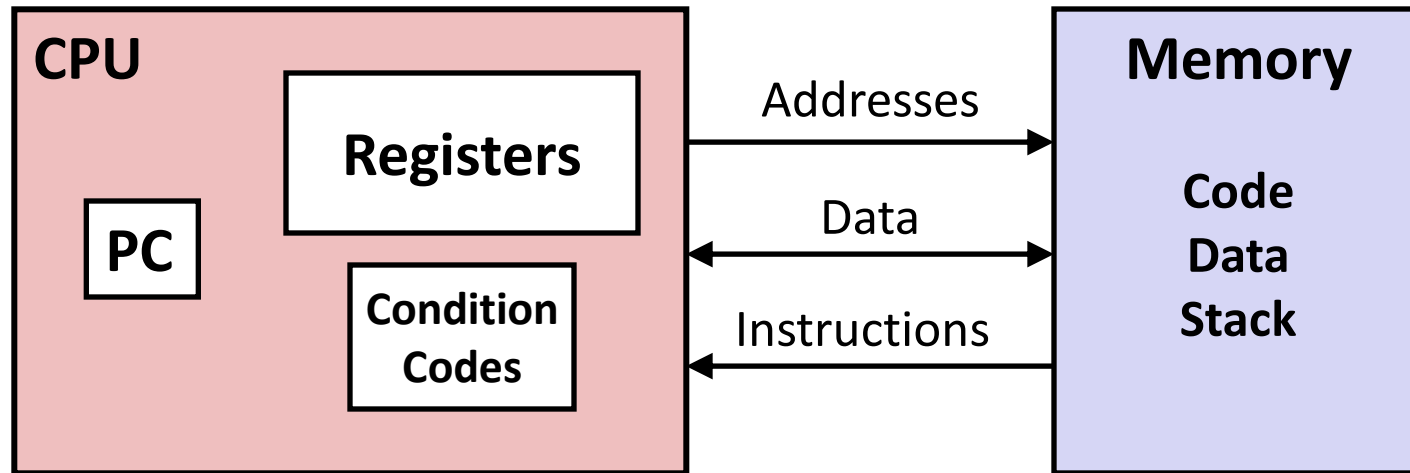
byoungyoung@snu.ac.kr

<https://lifeasageek.github.io>

Today

- From last week: Turning C into machine code
- Basics of control flow
- Condition codes
- Conditional branches
- Loops

Recall: ISA = Assembly/Machine Code View



Programmer-Visible State

- **PC: Program counter**
 - Address of next instruction
- **Register file**
 - Heavily used program data
- **Condition codes**
 - Store status information about most recent arithmetic or logical operation
 - Used for conditional branching

■ Memory

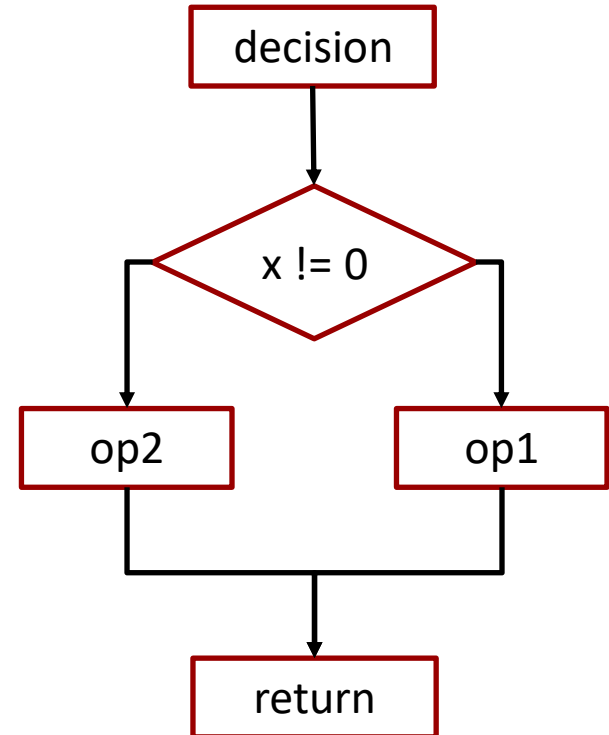
- Byte addressable array
- Code and user data
- Stack to support procedures



Control flow

```
extern void op1(void);  
extern void op2(void);
```

```
void decision(int x) {  
    if (x) {  
        op1();  
    } else {  
        op2();  
    }  
}
```



Control flow in assembly language

```
extern void op1(void);
extern void op2(void);

void decision(int x) {
    if (x) {
        op1();
    } else {
        op2();
    }
}
```

```
decision:
    subq    $8, %rsp
    testl   %edi, %edi
    je      .L2
    call    op1
    jmp     .L1
.L2:
    call    op2
.L1:
    addq    $8, %rsp
    ret
```

Processor State (x86-64)

■ Information about currently executing program

- Temporary data
(`%rax`, ...)
- Location of runtime stack
(`%rsp`)
- Location of current code control point
(`%rip`, ...)
- Status of recent tests
(`CF`, `ZF`, `SF`, `OF`)

Current stack top

Registers

<code>%rax</code>	<code>%r8</code>
<code>%rbx</code>	<code>%r9</code>
<code>%rcx</code>	<code>%r10</code>
<code>%rdx</code>	<code>%r11</code>
<code>%rsi</code>	<code>%r12</code>
<code>%rdi</code>	<code>%r13</code>
<code>%rsp</code>	<code>%r14</code>
<code>%rbp</code>	<code>%r15</code>

`%rip`

Instruction pointer

`CF`

`ZF`

`SF`

`OF`

Condition codes

Condition Codes (Implicit Setting)

■ Single bit registers

- **CF** Carry Flag (for unsigned) **SF** Sign Flag (for signed)
- **ZF** Zero Flag **OF** Overflow Flag (for signed)

■ Implicitly set (as side effect) after arithmetic operations

Example: `addq Src, Dest` $\leftrightarrow$ `t = a+b`

CF set if carry/borrow out from most significant bit (unsigned overflow)

ZF set if `t == 0`

SF set if `t < 0` (as signed)

OF set if two's-complement (signed) overflow

`(a>0 && b>0 && t<0) || (a<0 && b<0 && t>=0)`

ZF set when

000000000000...000000000000

SF set when

$$\begin{array}{r} \boxed{\text{yxxxxxxxxxxxxxxxxx} \dots} \\ + \boxed{\text{yxxxxxxxxxxxxxxxxx} \dots} \\ \hline \boxed{\text{1xxxxxxxxxxxxxxxxx} \dots} \end{array}$$

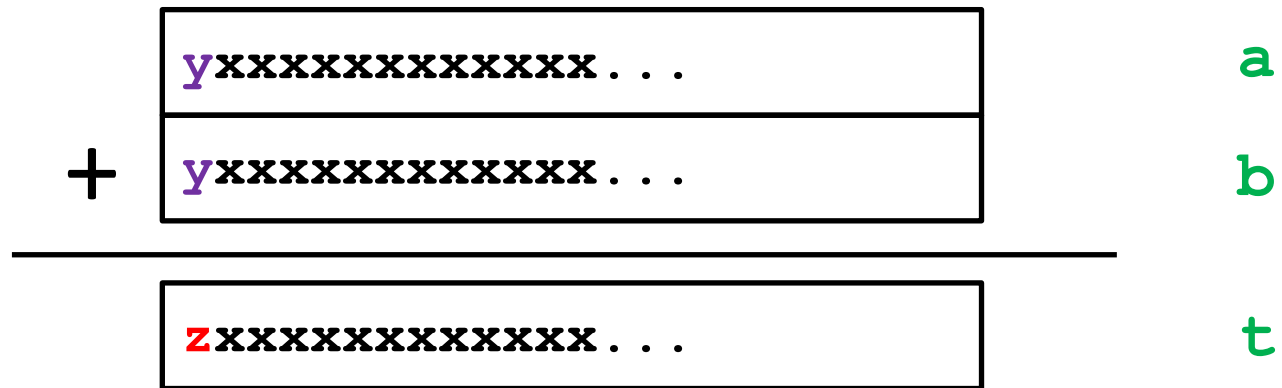
For signed arithmetic, this reports when result is a negative number

CF set when



For unsigned arithmetic, this reports overflow

OF set when



$$z = \sim y$$

`(a>0 && b>0 && t<0) || (a<0 && b<0 && t>=0)`

For signed arithmetic, this reports overflow

Condition Codes (Explicit Setting: Compare)

■ Explicit Setting by Compare Instruction

- `cmpq Src2, Src1`
- `cmpq b, a` like computing `a-b` without setting destination
- **CF set** if carry/borrow out from most significant bit
(used for unsigned comparisons)
- **ZF set** if `a == b`
- **SF set** if `(a-b) < 0` (as signed)
- **OF set** if two's-complement (signed) overflow
`(a>0 && b<0 && (a-b)<0) || (a<0 && b>0 && (a-b)>0)`

Condition Codes (Explicit Setting: Test)

■ Explicit Setting by Test instruction

- `testq Src2, Src1`
 - `testq b, a` like computing `a&b` without setting destination
- Sets condition codes based on value of `Src1` & `Src2`
- Useful to have one of the operands be a mask
- **ZF set** when `a&b == 0`
- **SF set** when `a&b < 0`

Very often:

```
testq    %rax, %rax
```

Today

- Control: Condition codes
- **Conditional branches**
- Loops

Jumping

■ jX Instructions

- Jump to different part of code depending on condition codes
- Implicit reading of condition codes

jX	Condition	Description
jmp	1	Unconditional
je	ZF	Equal / Zero
jne	$\sim ZF$	Not Equal / Not Zero
js	SF	Negative
jns	$\sim SF$	Nonnegative
jg	$\sim (SF \wedge OF) \ \& \ \sim ZF$	Greater (signed)
jge	$\sim (SF \wedge OF)$	Greater or Equal (signed)
jl	$SF \wedge OF$	Less (signed)
jle	$(SF \wedge OF) \ \ ZF$	Less or Equal (signed)
ja	$\sim CF \ \& \ \sim ZF$	Above (unsigned)
jb	CF	Below (unsigned)

Conditional Branch Example

■ Generation

```
$ gcc -Og -S -fno-if-conversion control.c
```

```
long absdiff
(long x, long y)
{
    long result;
    if (x > y)
        result = x-y;
    else
        result = y-x;
    return result;
}
```

```
absdiff:
    cmpq    %rsi, %rdi    # x-y
    jle     .L4
    movq    %rdi, %rax
    subq    %rsi, %rax    # x-y
    ret
.L4:       # x <= y
    movq    %rsi, %rax
    subq    %rdi, %rax    # y-x
    ret
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rax	Return value

Today

- Control: Condition codes
- Conditional branches
- **Loops**

“Do-While” Loop Example

C Code

```
long pcount_do(ulong x) {  
    long result = 0;  
    do {  
        result += x & 0x1;  
        x >>= 1;  
    } while (x);  
    return result;  
}
```

Goto Version

```
long pcount_goto(ulong x) {  
    long result = 0;  
loop:  
    result += x & 0x1;  
    x >>= 1;  
    if(x) goto loop;  
    return result;  
}
```

- Count number of 1's in argument *x*
- Use conditional branch to either continue looping or to exit loop

General “Do-While” Translation

C Code

```
do  
    Body  
while ( Test );
```

Goto Version

```
loop:  
    Body  
    if ( Test )  
        goto loop
```

“Do-While” Loop Compilation

```
long pcount_goto(ulong x) {  
    long result = 0;  
loop:  
    result += x & 0x1;  
    x >>= 1;  
    if(x) goto loop;  
    return result;  
}
```

Register	Use(s)
%rdi	Argument x
%rax	result

```
        xorq    %rax,%rax    #   result = 0  
.L2:    # loop:  
        movq    %rdi, %rdx  
        andl    $1, %edx     #   t = x & 0x1  
        addq    %rdx, %rax    #   result += t  
        shrq    %rdi         #   x >>= 1  
        jne     .L2          #   if(x) goto loop  
        ret
```

General “While” Translation

- “Jump-to-middle” translation
- Used with -Og

While version

```
while (Test)  
    Body
```



Goto Version

```
    goto test;  
loop:  
    Body  
test:  
    if (Test)  
        goto loop;  
done:
```

While Loop Example

C Code

```
long pcount_while(ulong x)
{
    long result = 0;
    while (x) {
        result += x & 0x1;
        x >>= 1;
    }
    return result;
}
```

Jump to Middle

```
long pcount_goto_jtm(ulong x)
{
    long result = 0;
    goto test;
loop:
    result += x & 0x1;
    x >>= 1;
test:
    if(x) goto loop;
    return result;
}
```

- Compare to do-while version of function
- Initial goto starts loop at test

“For” Loop Form

General Form

```
for (Init; Test; Update)  
  
    Body
```

```
#define WSIZE 8*sizeof(int)  
long pcount_for(ulong x)  
{  
    size_t i;  
    long result = 0;  
    for (i = 0; i < WSIZE; i++)  
    {  
        unsigned bit =  
            (x >> i) & 0x1;  
        result += bit;  
    }  
    return result;  
}
```

Init

```
i = 0
```

Test

```
i < WSIZE
```

Update

```
i++
```

Body

```
{  
    unsigned bit =  
        (x >> i) & 0x1;  
    result += bit;  
}
```

“For” Loop → While Loop

For Version

```
for (Init; Test; Update)  
    Body
```



While Version

```
Init;  
while (Test) {  
    Body  
    Update;  
}
```

Summarizing

■ C Control

- if-then-else
- do-while
- while, for
- Switch (didn't cover here)

■ Assembler Control

- Conditional jump
- Conditional move
- Indirect jump (via jump tables)
- Compiler generates code sequence to implement more complex control

Summary

■ Today

- Control: Condition codes
- Conditional branches & conditional moves
- Loops
- Switch statements

■ Next Time

- Stack
- Call / return
- Procedure call discipline